

Credo quia absurdum (?)

Proof Generation for Saturating First-Order Theorem Provers

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Agenda

Structure and Representation of Proofs

Proof Generation

Proof Applications

Conclusion

Structure and Representation of Proofs

Refutational Theorem Proving $\{A_1, A_2, \dots, A_n\} \models C$







Refutational Theorem Proving $\{A_1, A_2, \dots, A_n\} \models C$ iff

 $\{A_1, A_2, \ldots, A_n, \neg C\}$ is unsatisfiable







Refutational Theorem Proving $\{A_1, A_2, \ldots, A_n\} \models C$ iff $\{A_1, A_2, \ldots, A_n, \neg C\}$ is unsatisfiable Œ iff $cnf(\{A_1, A_2, A_n, \neg C\})$ is unsatisfiable

the future of theorem provin



Refutational Theorem Proving $\{A_1, A_2, \ldots, A_n\} \models C$ iff $\{A_1, A_2, \ldots, A_n, \neg C\}$ is unsatisfiable E iff $cnf(\{A_1, A_2, A_n, \neg C\})$ is unsatisfiable the future of theorem provin iff

$$cnf(\{A_1, A_2, A_n, \neg C\}) \vdash \Box$$



Ideal: Proofs as Sequences of Proof Steps

A derivation is a list of steps

Each step carries a clause/formula

Each step is either...

- Assumed (e.g. axioms, conjecture)
- Logically derived from earlier steps

A proof is a derivation that either...

- derives the conjecture
- derives a contradiction from the negated conjecture

Good mental model!

Reality: Proofs as Sequences of Proof Steps

Initial clauses/formulas

- Axioms/Conjectures/Hypotheses
- Justified by assumption

Derived clauses/formulas

- Justified by reference to (topologically) preceding steps
- Defined logical relationship to predecessors
 - Most frequent case: theorem of predecessors
 - Exceptions: Skolemization, negation of conjecture, ...

(Introduced definitions)

- Don't affect satisfiability/provability
- Justified by definition

Logical Languages for FOF

 Problems Otter "lists" LOP (CNF only) 	 Proofs/Derivations Otter "proof object" PCL (UEQ only)
 DFG TPTP (v1, v2) 	 DFG (but nobody uses DFG)

Logical Languages for FOF



Modern Convergence: TPTP v3

TPTP v3 language

Consistent syntax for different classes

- CNF is sub-case of FOF
- FOF is sub-case of TFF

Applicable for a wide range of applications

- Problem specifications
- Proofs/derivations
- Models

Easily parsable

- Prolog-parsable
- Lex/Yacc grammar
- Recursive-descent with 1-token look-ahead
- Widely used and supported
 - CASC
 - Major provers (E, SPASS, Vampire, iProver, ...)
 - Used by integrators

Example

```
fof(c 0 0, conjecture, (?[X3]:(human(X3)&X3!=john)), file('humen.p', someone not john)).
fof(c 0 1, axiom, (?[X3]:(human(X3)&grade(X3)=a)), file('humen.p', someone got an a)).
fof(c 0 2, axiom, (grade(john)=f), file('humen.p', john failed)).
fof(c_0_3, axiom, (a!=f), file('humen.p', distinct_grades)).
fof(c 0 4, negated conjecture, (~(?[X3]:(human(X3)&X3!=john))),
    inference(assume negation,[status(cth)],[c 0 0])).
fof(c_0_5, negated_conjecture, (![X4]:(~human(X4)|X4=john)),
    inference(variable rename, [status(thm)], [inference(fof nnf, [status(thm)], [c 0 4])])).
fof (c 0 6, plain, ((human(esk1 0)&grade(esk1 0)=a)),
    inference(skolemize,[status(esa)],[inference(variable_rename,[status(thm)],[c 0 1])])).
cnf(c 0 7, negated conjecture, (X1=john|~human(X1)),
    inference(split conjunct,[status(thm)],[c 0 5])).
cnf(c 0 8,plain,(human(esk1 0)),
    inference(split_conjunct,[status(thm)],[c_0_6])).
cnf(c \ 0 \ 9, plain, (grade(esk1 \ 0)=a),
    inference(split conjunct,[status(thm)],[c 0 6])).
cnf(c_0_10, negated_conjecture, (esk1_0=john),
    inference(spm,[status(thm)],[c 0 7, c 0 8])).
cnf(c 0 11,plain,(grade(john)=f),
    inference(split_conjunct,[status(thm)],[c_0_2])).
cnf(c 0 12, plain, (a!=f),
    inference(split conjunct,[status(thm)],[c 0 3])).
cnf(c 0 13, plain, ($false),
    inference(sr,[status(thm)],[inference(rw,[status(thm)],
    [inference(rw,[status(thm)],[c 0 9, c 0 10]), c 0 11]), c 0 12]), ['proof']).
```

























c_0_9: grade(esk1_0)=a c_0_10: esk1_0=john c_0_11:grade(john)=f c_0_12: a!=f







Compl[ie]mentary Example

TPTP v3 idiosyncrasies

No inference semantics

- Rules are just names
- Rules are system-dependent

Incomplete inference description

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- Syntactic support not widely supported

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Workarounds:

- Inference status
- Proof reconstruction
Proof Generation



Clausification and Saturation

Clausification

- Terminating
- (Usually) deterministic
- (Usually) non-destructive
- Sometimes done by external tool

Saturation

- Many degrees of freedom
- Arbitrary search time
- Generating inferences
 - Create new clauses
 - Necessary for completeness
- Simplifying inferences
 - Modify/remove existing clauses
 - Necessary for performance

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Recording clausification is straightforward

 ... but not always done

Efficiently recording saturation is difficult

 ...some settle for inefficient

Deduction vs. Simplification

Superposition
$$\frac{s \simeq t \lor S \quad u \neq v \lor R}{\sigma(u[p \leftarrow t] \neq v \lor S \lor R)}$$

if $\sigma = mgu(u|_{p}, s), [...]$
Rewriting
$$\frac{s \simeq t \quad u \neq v \lor R}{s \simeq t \quad u[p \leftarrow \sigma(t)] \neq v \lor R}$$

if $u|_{p} = \sigma(s)$ and $\sigma(s) > \sigma(t)$



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- Invariant: All generating inferences with premises from P have been performed
- Invariant: P is interreduced
- Clauses added to U are simplified with respect to P

Naive Proof Generation

Basic approach:

- Store (or dump) all intermediate proof steps
- Extract proof steps in post-processing

Problem: Necessary steps only known after the proof concludes

- Intermediate results are expensive to store
- Example: A ring with $X^4 = X$ is Abelian
 - Proof search (E): 5.4s
 - Proof search with inference dump: 11.4s
 - Post-processing: 17.6s
 - Temporary file size: 480 000 steps, 117MB
 - Proof size: 154 steps, 31 kB

Only suitable for small problems/short run-times

Optimized Proof Object Construction Observation: Only clauses in *P* are premises!



Optimized Proof Object Construction



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Proof recording:

- Simplified P-clauses are archived
- Clauses record their history
 - Inference rules
 - P-clauses involved

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 - ► Inference rules

 P-clauses involved

Proof extraction

- Track parent relation
- Topological sort
- Print proof

Optimized Proof Generation

Example: A ring with $X^4 = X$ is Abelian

Naive approach

- Proof search (E): 5.4s
- Proof search with inference dump: 11.4s
- Post-processing: 17.6s
- Temporary file size: 480 000 steps, 117MB
- Proof size: 154 steps, 31 kB
- Optimized approach
 - Proof search (E): 5.5s
 - Proof search with inference dump: -
 - Post-processing: -
 - Temporary file size: -
 - Proof size: 154 steps, 31 kB
- Example is typical
 - Optimized overhead: 0.24% over TPTP 5.4.0

Proof Applications

Why Proofs?

Trust

- in the proof
- in the ATP system
- in the specification
- Understanding
 - of the proof
 - of the domain
 - of the search process

Learning

- of important domain statements
- of search control information
- of the domain structure



Proof Checking

Semantic proof checking

- Step-by-step check
- Verify semantic status (conclusion can be derived "somehow" from premises)
- Use alternative theorem prover (or configuration)

Syntactic proof checking

- Show correctness of individual inference rule applications
- With TPTP syntax: Requires proof reconstruction
- E.g. Metis in Isabelle/Sledgehammer

Proof Visualization



Proof Visualization



Another Example



Another Example



(A ring with $X^4 = X$ is Abelian)

Interactive Visualization



Learning

Heuristics learning

- Find formulas that frequently appear in proofs
- Generalize and reuse

Axiom selection

. . .

 Learn relationship between conjecture and useful axioms



Conclusion

Efficient proof generation is non-trivial, but possible

TPTP v3 is a useful and used standard for proof representation

Proof objects are useful for trust building and learning

Use of proof objects is still in its infancy - we need more tools

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Proof presentation is a big open area

Ceterum Censeo...

Bug reports for E should include:

- The exact command line leading to the bug
- All input files needed to reproduce the bug
- A description of what seems wrong
- The output of eprover --version

